# **ESYNACKTIV**

Discovering and exploiting a kernel pool overflow on modern Windows 10

Sthack 2021

## ~:\$ whoami

- Fabien Perigaud
- @0xf4b on Twitter

#### Working for Synacktiv

- Offensive security
- 90 ninjas: pentest, reverse engineering, development
- We are hiring!

#### Reverse engineering team technical leader

- 30 reversers
- Reverse, vulnerability research and exploitation, low level dev



## **Agenda**

- This short introduction
- Pwn2Own Vancouver 2021
- Windows 10 Kernel Attack Surface
- Vulnerability Discovery
- Exploitation: "expecting shell root"
- Results and Conclusion

## Pwn2Own Vancouver 2021

### Pwn20wn



### "Hacking contest" by ZDI

Prove exploitation of devices or software in widespread use

#### Usually takes place twice a year

- Vancouver: virtualization, browsers, OS, Tesla, ...
- Tokyo: smartphones, "smart devices", routers, printers, ...
- Miami 2020: ICS/SCADA

### Pwn all the things!

- Payload should prove arbitrary code execution
- Remote/elevated shell, blinking leds on a router, image display on a printer...



### Pwn20wn - Rules



### Each category has its own rules...

- ... but rules can be adapted if kindly asked :)
- For example, enabling a non-default service...
- ... as long as this configuration matches reality!

### 5 minutes time slot, up to 3 attempts

- Exploit can be time consuming (race condition, huge allocations, ...)
- Stability can be perfectible :)



## Pwn20wn – Rules (2)



### For each target, only the first pwn is considered a WIN

- In practice, every successful pwn is a WIN...
- ... unless a vulnerability collision occurs...
- ... between contestants, or with the vendor!

#### Order of the contestants is drawn at random

- Small papers in a hat :)
- In the end, Synacktiv contestants usually get the last slot



## Pwn20wn - Rewards



# Each pwned target is rewarded, depending on the estimated difficulty

- Cash prize (way lower than the O-day market, but still interesting :))
- Master of Pwn points → additional cash prize for the "Master of Pwn"

#### Add-on bonus

- For some targets
- LPE, sandbox escape, ...
- Usually a few more points and dollars



### Pwn20wn Vancouver 2021

### Targets

- Desktop browsers (Chrome, Safari, Firefox, Edge)
- Enterprise applications (Office, Reader, Zoom, Teams)
- Server (RDP, Exchange, SharePoint)
- Automotive (Tesla Model 3)
- Local Privilege Escalation (Windows, Ubuntu)

### LPEs: "easier" targets!



## **Pwn20wn Vancouver 2021 — Windows 10 LPE**

Focus on the Windows 10 LPE

- Interesting execution context
  - Unprivileged user...
  - ... but no sandboxing!
    - Medium integrity level
- Vulnerability must be in the kernel



## Windows 10 Kernel Attack Surface

### **Kernel Attack Surface**



#### Ntoskrnl

- Windows kernel image
- Interrupts, memory management, kernel objects (processes, threads, files, registry, ...), syscalls and more
- Very interesting target, might be reachable from the hardest sandbox level
- Drawback: huge focus from security researchers

#### Win32k

- Huge graphic subsystem, own syscall table
- Old code base, many vulnerabilities
- Also reachable from some sandbox contexts
- Drawback: also a huge focus from security researchers



### **Kernel Attack Surface - Drivers**

#### Drivers

- PE loaded in Kernel-land
- ".sys" file on the disk
- Usually linked to a service

#### Userland access

- Driver create a Device object "XXX"
- Userland opens the device through "\\?\GLOBALROOT\Device\XXX"



### **Kernel Attack Surface - Devices**

### List all the devices present on a default Windows 10

- WinDBG to the rescue!
- !object \Device → 159 devices

```
0: kd> !object \Device
Object: ffff988952031060 Type: (ffffaf07a6c71900) Directory
    ObjectHeader: ffff988952031030 (new version)
    HandleCount: 2 PointerCount: 65717
    Directory Object: ffff988952041e60 Name: Device
```

Hash	Address	Type	Name
00	ffffaf07adaafcb0	Device	00000030
	ffffaf07ad94a050	Device	NDMP2
	ffffaf07a82c8360	Device	NTPNP_PCI0002
01	ffffaf07ad9d2050	Device	NDMP3

. . .

### **Check devices access**



### Dumb/dirty way to check access rights

- Try to open device with R/W access
- Fast to write with Python ctypes

```
for dev in devices:
    file_handle = windll.kernel32.CreateFileA("\\\?\\GLOBALROOT\\Device\\
%s" % dev, GENERIC_READ | GENERIC_WRITE, 0, NULL, OPEN_EXISTING,
FILE_ATTRIBUTE_NORMAL, 0)
    if file_handle != INVALID_HANDLE_VALUE:
        print("[*] OK %s" % dev)
```



## Check devices access (2)

```
C:\Users\unpriv\Desktop>check devices.py
[*] OK Afd
[*] OK ahcache
[*] OK Beep
[*] OK CNG
[*] OK gpuenergydrv
[*] OK KsecDD
[*] OK LanmanDatagramReceiver
[*] OK Mailslot
[*] OK Mup
[*] OK NamedPipe
[\ldots]
[*] OK Netbios
[*] OK Null
[*] OK PEAuth
[*] OK RdpBus
[*] OK Spaceport
[*] OK Tdx
[*] OK USBFD0-0
[*] OK USBPD0-0
[*] OK VBoxGuest
[*] OK WindowsTrustedRT
                                        SYNACKTIV
[*] OK WINDRVR6
```

Afd	ahcache	Веер	CNG	gpuenergydrv
KsecDD	LanmanDatagramReceiver	Mailslot	Mup	NamedPipe
Netbios	Null	PEAuth	RdpBus	Spaceport
Tdx	USBFDO-0	USBPDO-0	VBoxGuest	
WinDRV6	WindowsTrustedRT			

Avd	ahcache	Beep	CivG	gpuenergydrv
KsecDD	LanmanDatagramReceiver	Marislot	Mup	NamedPipe
Netrios	Parall	PEAuth	RdpBus	Spaceport
Totx	USBFDO-0	USBPDO-0	VBoxSuest	
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Ard	ahcache	Beep	CivG	gpuenergydrv
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Tetx	USB DO-0	USBFDO-0	VBoxSuest	
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Ard	aheache	Beep	CivG	gpuenergydrv
KsecDD	LanmanDatagramReceiver	Marislot	Mup	NamedPipe
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Tox	USB DO-0	USBFDO-0	VBoxouest	
WinDRV6	WindowsTrustedRT			

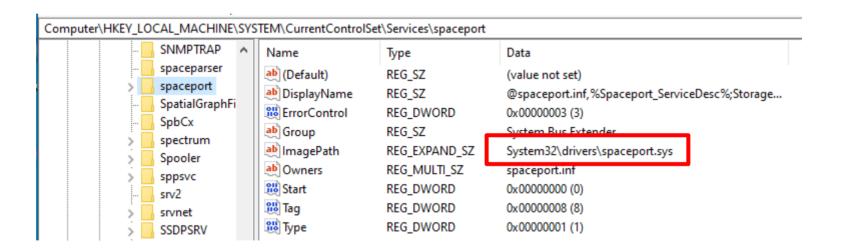
## Find the corresponding driver

### WinDBG again



## Find the corresponding driver

Corresponding service in the registry





# Vulnerability Discovery

### **Driver interaction**



### Driver Object can have 28 defined Major Functions

- Open / Close
- Read / Write
- IOCTL
- Etc.

### Usually defined in the DriverEntry function

- IOCTLs are usually the first thing to look at ...
- ... but others are also less analyzed!



## **Driver interaction — Major Functions**

```
DriverObject->MajorFunction[0] = SpSuccess;
DriverObject->MajorFunction[2] = SpSuccess;
DriverObject->MajorFunction[3] = (PDRIVER DISPATCH)SpDispatch;// IRP MJ READ
DriverObject->MajorFunction[4] = (PDRIVER DISPATCH)SpDispatch;// IRP MJ WRITE
DriverObject->MajorFunction[9] = (PDRIVER DISPATCH)SpDispatch;// IRP MJ FLUSH BUFFERS
DriverObject->MajorFunction[13] = (PDRIVER DISPATCH)SpDispatch;// IRP MJ FILE SYSTEM CONTROL
DriverObject->MajorFunction[14] = (PDRIVER DISPATCH)SpDispatch;// IRP MJ DEVICE CONTROL
DriverObject->MajorFunction[15] = (PDRIVER DISPATCH)SpDispatch;// IRP MJ INTERNAL DEVICE CONTROL
DriverObject->MajorFunction[16] = (PDRIVER DISPATCH)SpDispatch;// IRP MJ SHUTDOWN
DriverObject->MajorFunction[18] = SpSuccess;
DriverObject->MajorFunction[22] = (PDRIVER DISPATCH)SpDispatch;// IRP MJ POWER
DriverObject->MajorFunction[23] = (PDRIVER DISPATCH)SpDispatch;// IRP MJ SYSTEM CONTROL
DriverObject->MajorFunction[27] = (PDRIVER DISPATCH)SpDispatch;// IRP MJ PNP
*(&SpControlDispatchTable + 14) = (int ( stdcall * near *)(struct DEVICE OBJECT *, struct IRP *))SpControlDeviceControl;
*(&SpControlDispatchTable + 15) = (int ( stdcall * near *)(struct DEVICE OBJECT *, struct IRP *))SpControlScsi;
*(&SpControlDispatchTable + 16) = (int ( stdcall * near *)(struct DEVICE OBJECT *, struct IRP *))SpControlShutdown;
*(&SpControlDispatchTable + 27) = (int ( stdcall * near *)(struct DEVICE OBJECT *, struct IRP *))SpControlPnp;
```



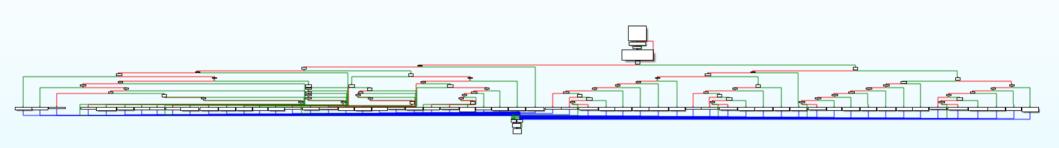
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*(&SpControlDispatchTable + 14) = (int ( stdcall * near *)(struct DEVICE OBJECT *, struct IRP *))SpControlDeviceControl
*(&SpControlDispatchTable + 15) = (int ( stdcall * near *)(struct DEVICE OBJECT *, struct IRP *))SpControlScs1;
*(&SpControlDispatchTable + 16) = (int ( stdcall * near *)(struct DEVICE OBJECT *, struct IRP *))SpControlShutdown;
*(&SpControlDispatchTable + 27) = (int ( stdcall * near *)(struct DEVICE OBJECT *, struct IRP *))SpControlPnp;
```



## **SpControlDeviceControl**

#### 58 handled IOCTLs



#### Manual review

- Focus on "Sploctl[Create|Set].\*"
- Ignore when privileges are checked ("SpAccessCheck.\*")



## **SpControlDeviceControl**

### Many checks are missing

When the research was performed :)

#### Several vulnerabilities have been found!

- Constraints to be reachable
- Might require creation of objects from a privileged context
  - Might be OK in real life, not for the contest

### One really interesting vulnerability



- No permission check
- Buffer overflow 101
  - Controlled memcpy size
  - Controlled content

```
v8->sp_list.Flink = &v7->sp_list;
v7->sp_list.Blink = &v8->sp_list;
p_sp_work->sp_list.Blink = &p_sp_work->sp_list;
p_sp_work->sp_list.Flink = &p_sp_work->sp_list;
p_sp_work->work_info->error_code = user_buffer->error_code;
p_sp_work->work_info->content_size = user_buffer->content_size;
memmove(
    (char *)p_sp_work->work_info + (unsigned int)p_sp_work->work_info->content_offset,
    (char *)user_buffer + (unsigned int)user_buffer->content_offset,
    (unsigned int)user_buffer->content_size);
KeSetEvent((PRKEVENT)&p_sp_work->set_event, 0, 0);
```



- No permission check
- Buffer overflow 101
  - Controlled memcpy size
  - Controlled content

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v7->sp_list.Blink = &v8->sp_list;
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p_sp_work->sp_list.Flink = &p_sp_work->sp_list;
p_sp_work->work_info->error_code = user_buffer->error_code;
p_sp_work->work_info->content_size = user_buffer->content_size;
memmove(
    (char *)p sp work->work info + (unsigned int)p sp work->work info->content offset,
    (char *)user_buffer + (unsigned int)user_buffer->content_offset,
    (unsigned int)user_buffer->content_size);
KeSetEvent((PRKEVENT)&p sp work->set event, 0, 0);
```

Dest: pool (heap) buffer



- No permission check
- Buffer overflow 101
  - Controlled memcpy size
  - Controlled content

```
v8->sp_list.Flink = &v7->sp_list;
v7->sp_list.Blink = &v8->sp_list;
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```

Source: user controlled buffer



- No permission check
- Buffer overflow 101
  - Controlled memcpy size
  - Controlled content

```
v8->sp_list.Flink = &v7->sp_list;
v7->sp_list.Blink = &v8->sp_list;
p_sp_work->sp_list.Blink = &p_sp_work->sp_list;
p_sp_work->sp_list.Flink = &p_sp_work->sp_list;
p_sp_work->work_info->error_code = user_buffer->error_code;
p_sp_work->work_info->content_size = user_buffer->content_size;
memmove(
    (char *)p_sp_work->work_info + (unsigned int)p_sp_work->work_info->content_offset,
    (char *)user_buffer + (unsigned int)user_buffer->content_offset,
    (unsigned int)user_buffer->content_size);
KeSetEvent((PRKEVENT)&p_sp_work->set_event, 0, 0);
```

Size: from user controlled buffer



## SploctlSetControlWork — Reaching the bug

- SploctlSetControlWork looks for a SP\_WORK\_INFO in a doubly-linked list ("LIST1")
  - Identified by a provided ID
- "LIST1" is populated by another IOCTL: SploctIGetControlWork
  - Gets an entry from another doubly-linked list ("LIST2")
  - Put it in "LIST1" and return its ID
- "LIST2" is populated by SP\_CONTROL\_WORK::Run
  - Reachable from several IOCTLs
  - SploctlAttachSpaceRemote is a good candidate

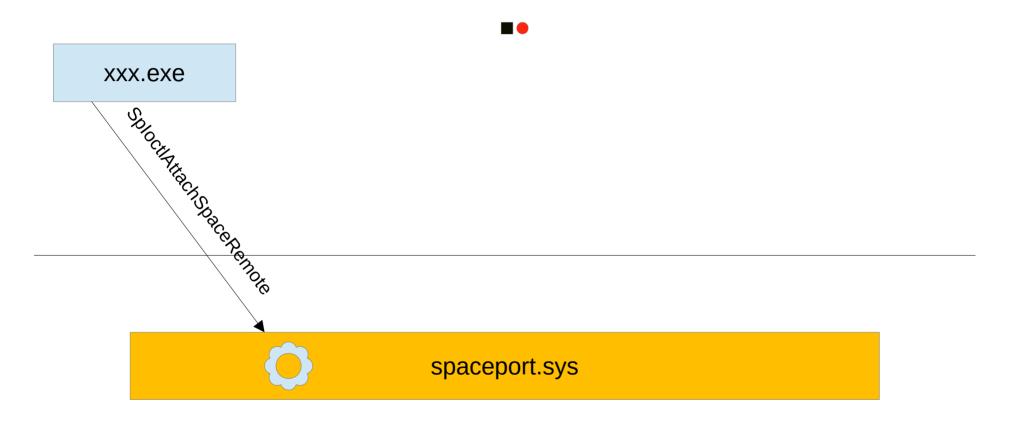


## **Workflow**

- If we call the IOCTLs sequentially
  - Stuck when calling SploctlGetControlWork → we never get an ID
- What's happening?

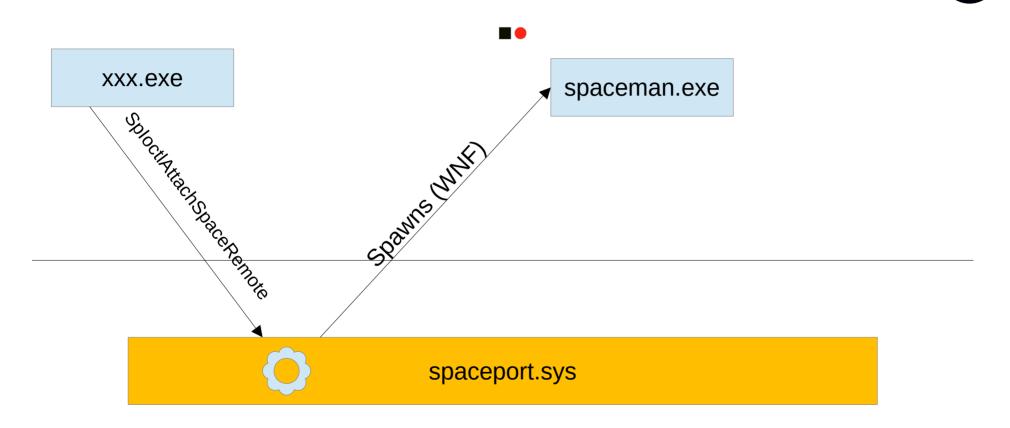


## **Usual workflow**



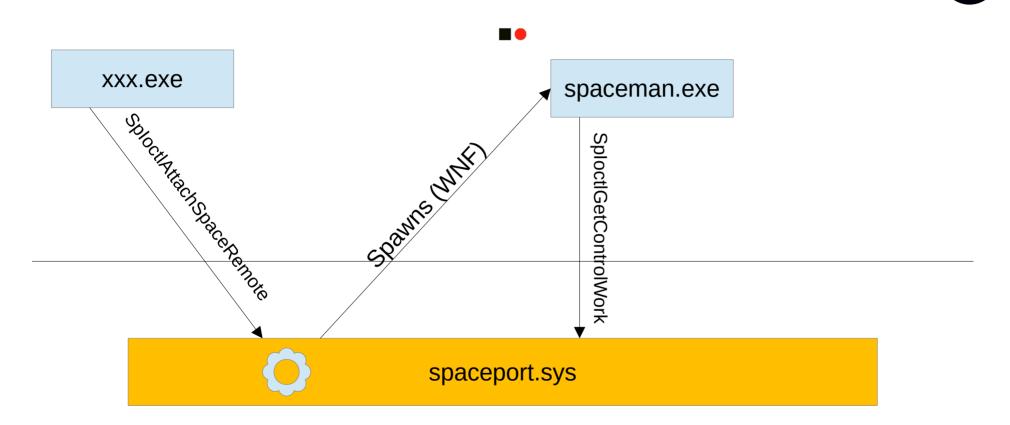


## **Usual workflow**



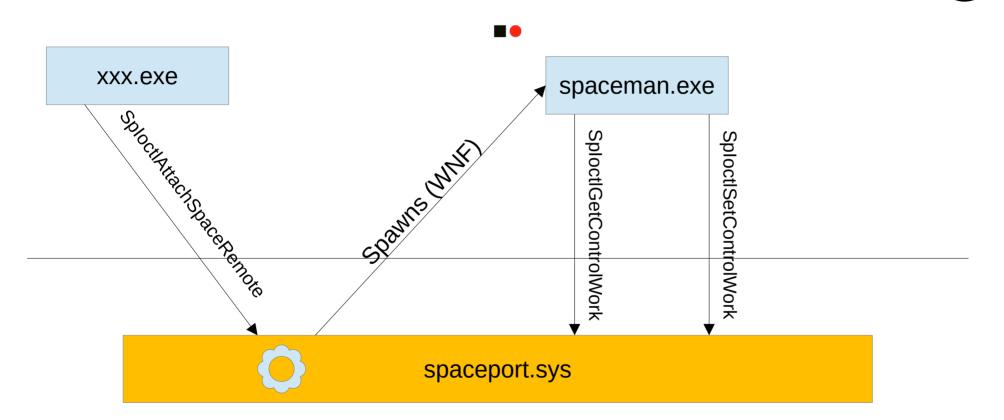


## **Usual workflow**





## **Usual workflow**





## **Usual Workflow**

- If we call the IOCTLs sequentially, we get raced by spaceman.exe
  - No more ID to be retrieved by SploctIGetControlWork

### However...

- ... meet asynchronous DeviceloControl!
- Call SploctIGetControlWork before SploctIAttachSpaceRemote
  - List is empty, driver puts the request on hold
  - IRP is queued, and dequeued when SP\_CONTROL\_WORK::Run is executed
- When spaceman.exe is executed, it has been raced by us :)



### Poc

- Thread #1 → call SploctlGetControlWork
  - Async, wait for result
- Thread #2 → call SploctlAttachRemoteSpace
  - Blocked until someone issues a SploctlSetControlWork
- Thread #1 → ID retrieved, call SploctlSetControlWork with bogus length



### PoC - BSOD

```
rax=ffffaf07b1584ed0 rbx=0000000000000000 rcx=ffffaf07f39a9112
rip=fffff80465dfa343 rsp=ffffc106ae2686d8 rbp=fffff80465e32048
iopl=0
          nv up ei ng nz na po nc
spaceport!memcpy+0x203:
fffff804`65dfa343 0f104411f0
                        movups xmm0,xmmword ptr [rcx+rdx-10h]
ds:ffffaf07`ee1c6292=?????????????????????????????????
ffffc106`ae2686d8 fffff804`65df6f19 : [...] : spaceport!memcpy+0x203
ffffc106`ae2686e0 fffff804`65e4c12b : [...] : spaceport!SP CONTROL WORK::Set+0xd9
ffffc106`ae268740 fffff804`65e41d7e : [...] : spaceport!SpIoctlSetControlWork+0x5b
ffffc106`ae268780 fffff804`65df3f50 : [...] : spaceport!SpControlDeviceControl+0x2ee
```



# Exploitation: "expecting shell root"

# **Vulnerability primitive**

#### Pool overflow

The pool is the Windows Kernel heap

### Target allocation

- SP\_WORK\_INFO allocation is made in the NonPagedPoolNx
- Size is 0x160 bytes
- Lies in the LFH (Low Fragmentation Heap)

#### Overflow constraints

- None :)
- We control content and size



# **Mitigations**

### kASLR

- Not a problem, Medium integrity level
- Various kernel APIs to get objects and modules addresses

### DEP / SMEP / CFG

- Kernel code execution is hard
- Data-only exploitation ftw!

### SMAP?

Only in a few contexts, not in ours :)



# **Exploitation strategy**

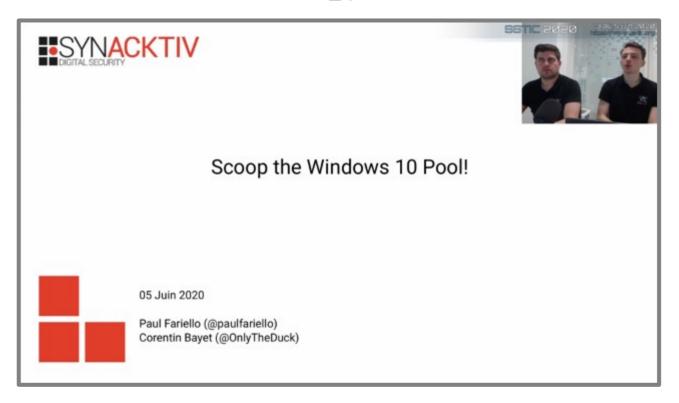
- Data only
- We want to run an elevated cmd.exe

- Target: process Token!
  - Swap token with a privileged one? (System)
  - Or enable powerful privileges!
- Let's turn our pool overflow into something interesting!



## La French Tech — SSTIC 2020







## La French Tech — SSTIC 2020



### Aligned Chunk Confusion

- New generic pool overflow exploitation method
- Abuses the CacheAligned bit in the POOL\_HEADER
- Read their paper for details!

### Requirements

Shape the pool to control the chunk after the vulnerable one



# **Pool massaging**

## Spray a bunch of 0x160 bytes allocations in NonPagedPoolNx

- We can use pipe objects
- PipeQueueEntry is in the NonPagedPoolNx, and we can control its size

| QueueEntry |
|------------|------------|------------|------------|------------|------------|------------|------------|------------|
| ,          | ,          | ,          | , ,        | ,          | ,          | , ,        | , ,        | ,          |



# **Pool massaging**

### Spray a bunch of 0x160 bytes allocations in NonPagedPoolNx

- We can use pipe objects
- PipeQueueEntry is in the NonPagedPoolNx, and we can control its size

QueueEntry QueueEntry QueueEntry QueueEntry QueueEntry QueueEntry QueueEntry QueueEntry QueueEntry
----------------------------------------------------------------------------------------------------

### Free a few ones to create holes

QueueEntry	Free	QueueEntry	QueueEntry	QueueEntry	QueueEntry	Free	QueueEntry	QueueEntry
------------	------	------------	------------	------------	------------	------	------------	------------



# **Pool massaging**

## Spray a bunch of 0x160 bytes allocations in NonPagedPoolNx

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QueueEntry   Queue	QueueEntry Queu	ueueEntry (	QueueEntry						
--------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------	-----------------	-------------	------------	------------	------------	------------	------------	------------	------------

### Free a few ones to create holes

QueueEntry	Free	QueueEntry	QueueEntry	QueueEntry	QueueEntry	Free	QueueEntry	QueueEntry
_		_		_	_		_	•

### Our vuln chunk should lie in one of these holes

QueueEntry	Free	QueueEntry	QueueEntry	QueueEntry	QueueEntry	Vuln Chunk	QueueEntry	QueueEntry
_		,	,	, ,	•		, .	, ,

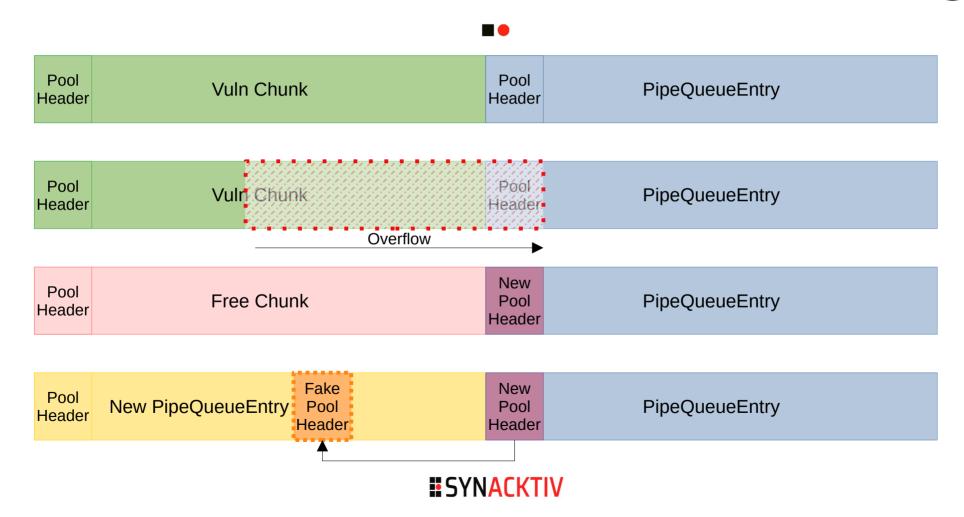


# Creating an overlapping chunk

- We can now trigger the overflow to change the next chunk (called victim) POOL\_HEADER
  - poolType set to 4 → CacheAligned
  - previousSize set to 0x100
- When freeing the victim, the allocator will look for a second POOL\_HEADER 0x100 bytes before the chunk
  - This creates a fake chunk of 0x260 bytes (0x160+0x100)
- Caveat: after exploitation, our vuln chunk is freed
  - We reuse it with a controlled PipeQueueEntry!



# **Creating an overlapping chunk - Graphics**

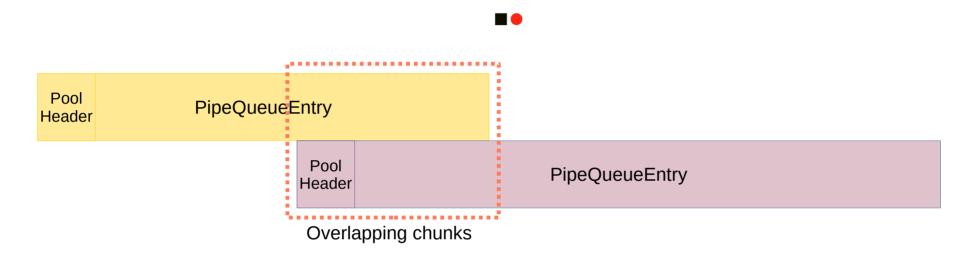


# Creating an overlapping chunk — Next steps

- We chose a size of 0x260 bytes for our fake chunk
  - LFH is for allocations < 0x200
  - Bigger allocations are handled by the VS (variable size) allocator
  - Lookaside lists can be enabled for faster allocations!
- When the victim chunk is freed
  - Fake POOL\_HEADER is read
  - Fake chunk is added to the 0x260 bytes lookaside list
- We can now make a new allocation of 0x260 bytes to reuse the fake chunk!



# Creating an overlapping chunk — Result



### Leak by reading the first pipe

Gives the PipeQueueEntry structure content



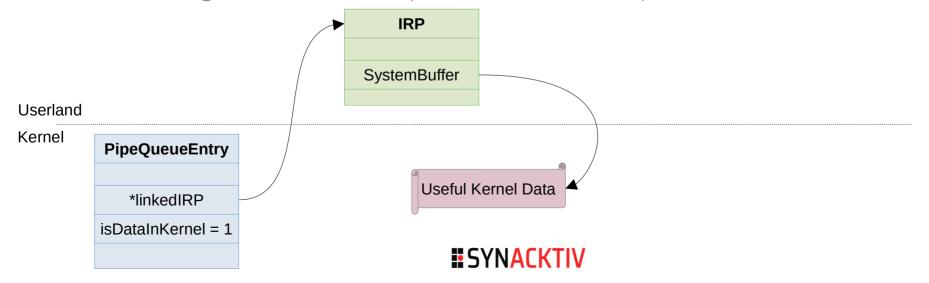
# **PipeQueueEntry structure**

```
List of PipeQueueEntry structures
struct PipeOueueEntry {
   LIST ENTRY list;
   IRP *IRP:
   uint64 t security;
   int isDataInKernel;
   int remaining bytes;
   int DataSize;
   int field 2C;
                           if (PipeQueueEntry->isDataInKernel == 1)
   char data[0];
                              data ptr = (PipeQueueEntry->linkedIRP->SystemBuffer);
};
                           else
                              data ptr = PipeQueueEntry->data;
```

### **#SYNACKTIV**

# **Arbitrary Read Primitive**

- Free the first PipeQueueEntry
- Reuse the chunk
  - We can overwrite the second PipeQueueEntry structure
  - Change the linkedIRP pointer to make it point to userland



# **Attacking the ProcessBilled pointer**

### POOL\_HEADER has a ProcessBilled field

- Obfuscated pointer to an EPROCESS
- If PoolQuota flag is set, EPROCESS→QuotaBlockPtr→value is decremented when allocation is freed

## Arbitrary decrement primitive

- Requires ability to forge a new obfuscated pointer
- ProcessBilled == @EPROCESS ^ @chunk ^ ExpPoolQuotaCookie



# Finding ExpPoolQuotaCookie

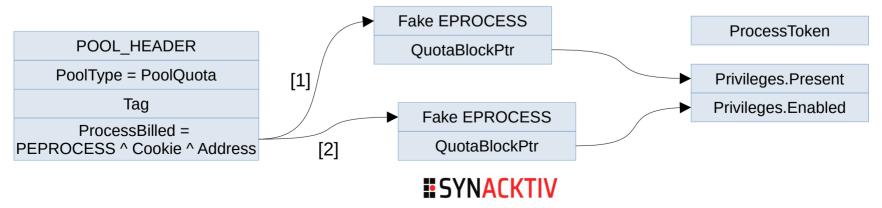


- From Medium Integrity Level, kernel APIs can be used to get a kernel address from a handle
  - NtQuerySystemInformation(SystemHandleInformation)
  - We have our EPROCESS address!
- In the previous leak
  - We got our PipeQueueEntry address through the doubly-linked list
  - We got its POOL\_HEADER ProcessBilled
- ExpPoolQuotaCookie can be computed



# Arbitrary decrement to privilege escalation

- A process Token contains its privileges
- Strategy: decrement "Enabled" and "Present" fields to enable SeDebugPrivilege
  - We build a fake EPROCESS which QuotaBlockPtr points to the target field
- 2 more reuses needed to change the ProcessBilled twice



# **SeDebugPrivileges**

Allow debugging every process on the system

### Strategy

- Open winlogon.exe
- Inject shellcode
- Spawn a SYSTEM cmd.exe

### Quick demo!



# Results and Conclusion

## **Pwn20wn results**

Exploit worked at first attempt \o/

- Debriefing with ZDI
  - Bug is unknown to ZDI \o/ \o/
- Debriefing with Microsoft
  - ... bug is already known to Microsoft ...
  - They proved it by showing the bug report
- Partial Win :(



## **Timeline**

- ?????? Vulnerability reported to Microsoft by "vbty" (according to Microsoft advisory)
- 8 April 2021 Vulnerability exploited during P2O
- 13 July 2021 Vulnerability fixed by Microsoft (CVE-2021-33751)

#### Exploitability

The following table provides an exploitability assessment for this vulnerability at the time of original publication.

Publicly Disclosed	Exploited	Exploitability Assessment
No	No	Exploitation Less Likely



## **Fix**

```
if ( user buffer->content size <= (unsigned int)work info->content size )
  FILINK = (SP WORK */p Sp WORK->Sp LIST.FILINK;
  if ( (SP WORK *)p sp work->sp list.Flink->Blink != p sp work
    [] (Blink = (SP WORK *)p sp work->sp list.Blink, (SP WORK *)Blink->sp list.Flink != p sp work) )
      fastfail(3u);
  Blink->sp list.Flink = &Flink->sp list;
  Flink->sp_list.Blink = &Blink->sp list;
  p sp work->sp list.Blink = &p sp work->sp list;
  p sp work->sp list.Flink = &p sp work->sp list;
  p sp work->work info->error code = user buffer->error code;
  p sp work->work info->content size = user buffer->content size;
  memmove(
    (char *)p_sp_work->work_info + (unsigned int)p_sp work->work info->content offset.
    (char *)user buffer + (unsigned int)user buffer->content offset,
    (unsigned int)user buffer->content size);
```



## Free Oday?

```
if ( user buffer->content size <= (unsigned int)work info->content size )
  Flink = (SP WORK *)p sp work->sp list.Flink;
  if ( (SP WORK *)p sp work->sp list.Flink->Blink != p sp work
    [] (Blink = (SP WORK *)p sp work->sp list.Blink, (SP WORK *)Blink->sp list.Flink != p sp work) )
    _fastfail(3u);
  Blink->sp list.Flink = &Flink->sp list;
  Flink->sp list.Blink = &Blink->sp list;
  p sp work->sp list.Blink = &p sp work->sp list;
  p sp work->sp list.Flink = &p sp work->sp list;
  p sp work->work info->error code = user buffer->error code;
  p sp work->work info->content size = user buffer->content size;
  memmove(
    (char *)p sp work->work info + (unsigned int)p sp work->work info->content offset,
    (char *)user buffer + (unsigned int)user buffer->content offset,
    (unsigned int)user buffer->content size);
```



## **Additional fix...**



## **Final words**

## A generic pool overflow exploitation method exists!

- ... and works on real cases!
- Thanks to the new kernel pool from 19H1

### Try Pwn2Own!

- Some targets do not require so much effort
- Attack surface is quite huge!





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